



Q.1) State whether following are Tautologies or contradiction?

- a) $\sim \{(\sim p) \wedge (\sim q)\} \Leftrightarrow (p \vee q)$;
 b) $\{p \wedge (p \Rightarrow q)\} \Rightarrow q$;
 c) $(p \vee q) \wedge \{p \vee (\sim q)\} \wedge \{(\sim p) \vee q\} \wedge \{(\sim p) \vee (\sim q)\}$
 d) $[(p \wedge q) \vee \{q \wedge (\sim r)\}] \Leftrightarrow [\{(\sim p) \wedge r\} \vee \{(\sim q) \wedge (\sim r)\}]$
 e) $\sim \{p \wedge (\sim q)\}$;
 f) $[(p \Rightarrow q) \wedge (q \Rightarrow r)] \Rightarrow (p \Rightarrow r)$

Q.2) Prove the following

- a) $(a + b).(a' + c) = (a.c) + (a'.b)$, $\forall a,b,c \in B$
 b) $a.b + a.b' + a'.b + a'.b' = 1$, $\forall a,b,c \in B$
 c) $p.q.r + p.q.r' + p.q'.r + p'.q.r = p.q + q.r + r.p$, $\forall p,q,r \in B$

Q.3) State & Prove De Morgan's Theorem ?

Q.4) Express the following functions into disjunctive normal form:

- a) $x.y$ b) $x + x'.y$ c) $f(x,y,z) = x.y' + x.z + x.y$ d) $f(x,y,z) = (x + y + z).(xy + x'z)'$
 e) $(x.y' + x.z)' + x'$ f) $(x + y)(x + y')(x' + z)$ g) $(x' + y)'.(x + z)' + (y.z)'$

Q.5) Express the following functions into conjunctive normal form:

- a) $f(x,y) = x + x'y$ b) $f(x,y,z) = xy' + xz + xy$ c) $(x + y + z)(xy + x'z)'$ d) $(x + y)(x + y')(x' + z)$

Answers

Ans.1 a) T b) T c) F d) F e) T f) T

Ans.4 a) $xyz + xyz'$ b) $xy + xy' + x'y$ c) $xyz + xy'z + xyz' + xy'z'$ d) $xy'z + xy'z' + x'yz'$
 e) $x'yz + x'yz' + x'y'z + x'y'z' + xyz$ f) $xyz + xy'z$ g) $xy'z + xy'z' + x'y'z + xyz' + x'yz' + x'y'z'$

Ans.5 a) $(x+y)$ b) $(x+y+z)(x+y'+z)(x+y+z')(x+y'+z')$
 c) $(x+y+z)(x'+y'+z)(x'+y'+z')(x+y+z')(x+y'+z')$ d) $(x+y+z)(x+y+z')(x+y'+z)(x+y'+z')(x'+y+z)(x'+y'+z)$



Truth Table of Logical Connectives

p	q	$p \wedge q$	$p \vee q$	$p \rightarrow q$	$p \leftrightarrow q$
T	T	T	T	T	T
T	F	F	T	F	F
F	T	F	T	T	F
F	F	F	F	T	T

p	q	$p \vee q$
T	T	T
T	F	T
F	T	T
F	F	F

Some Important Boolean Algebraic Laws

1.	Law of Identity	$A = A$ $\overline{\overline{A}} = A$
2.	Commutative Law	$A \cdot B = B \cdot A$ $A + B = B + A$
3.	Associative Law	$A \cdot (B \cdot C) = A \cdot B \cdot C$ $A + (B + C) = A + B + C$
4.	Idempotent Law	$A \cdot A = A$ $A + A = A$
5.	Double Negative Law	$\overline{\overline{A}} = A$
6.	Complementary Law	$A \cdot \overline{A} = 0$ $A + \overline{A} = 1$
7.	Law of Intersection	$A \cdot 1 = A$ $A \cdot 0 = 0$
8.	Law of Union	$A + 1 = 1$ $A + 0 = A$
9.	DeMorgan's Theorem	$\overline{AB} = \overline{A} + \overline{B}$ $\overline{A + B} = \overline{A} \overline{B}$
10.	Distributive Law	$A \cdot (B + C) = (A \cdot B) + (A \cdot C)$ $A + (BC) = (A + B) \cdot (A + C)$
11.	Law of Absorption	$A \cdot (A + B) = A$ $A + (AB) = A$
12.	Law of Common Identities	$A \cdot (\overline{A} + B) = AB$ $A + (\overline{A}B) = A + B$



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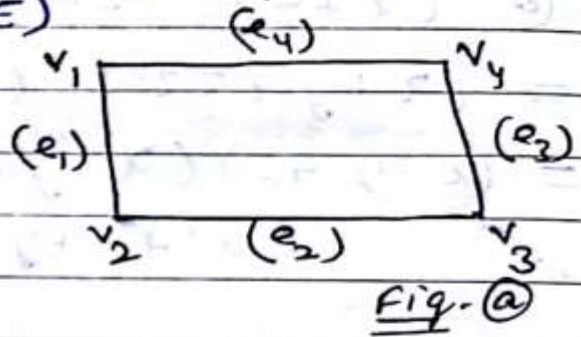
Graph:- A graph is set of objects whose elements or vertices i.e. nodes (V) and branches i.e. Edges (E)

we denote graph as $G = (V, E)$

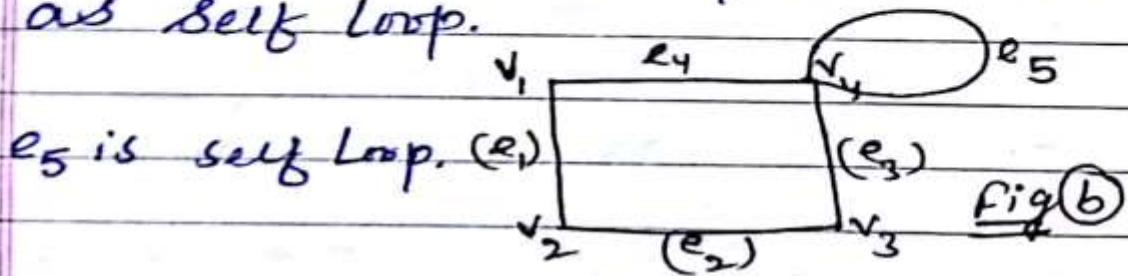
Here

$$E = \{e_1, e_2, e_3, e_4\}$$

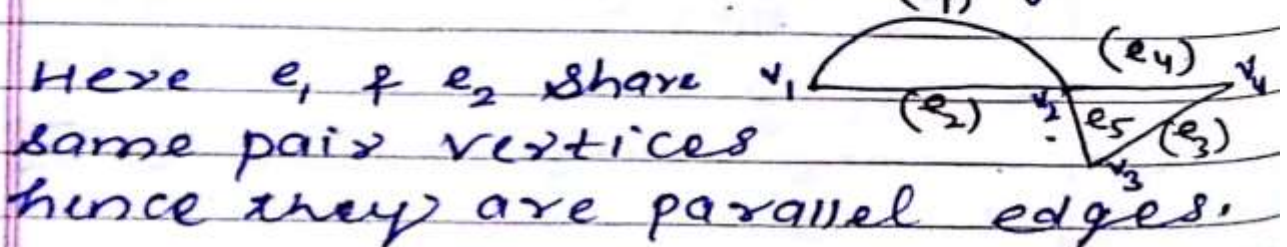
$$V = \{v_1, v_2, v_3, v_4\}$$



Self Loop:- If an edge starts and end on (LOOP) the same vertex, it is called as self loop.



Parallel Edges:- If two or more than two edges starts and end with the same pair of vertices, then they are said to be parallel edges.





Simple Graph: A graph that has neither self loop nor parallel edge is called a simple graph.

Incident edges: In fig. (b) e_3, e_4 & e_5 are incident on vertex v_4 .

Adjacent edge: Two nonparallel edges which incident on common vertex is called adjacent edge.
In fig. (b) e_3 & e_4 are adjacent edges.

Adjacent vertex: Two vertex are said to be an adjacent vertex, if there is an edge joining them.
In fig. (b), v_1 & v_2 are adjacent vertex.

Degree of vertex: No. of edges incident on a vertex, give the degree of the vertex.
 $\text{deg}(v_4) = 4$ in fig (b) where self loop is counted twice.
A vertex with odd degree is called as odd vertex whereas a vertex with even degree is called as even vertex.

Regular Graph: A graph with all vertices of same degree, is called regular graph fig (a)



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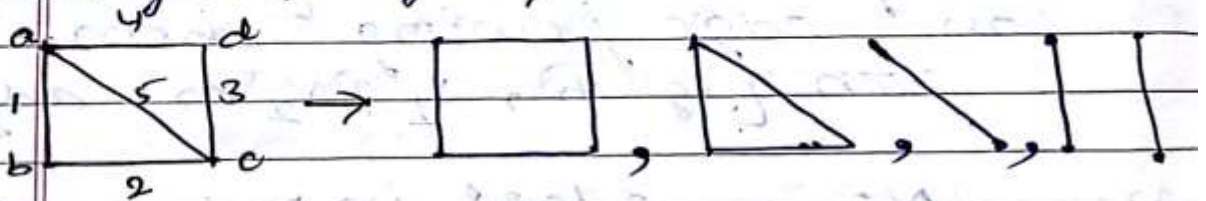
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Isolated vertex: A vertex of degree zero, is called isolated vertex or end vertex.

Pendant vertex: A vertex with degree 1 is called pendant vertex.

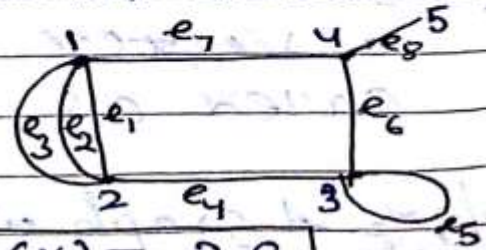
Null graph: A graph having no edges is called a null graph.

Subgraph: Subgraph of any graph is obtained by deletion of some vertices or edges from the original graph.



Theorem: The sum of the degrees of all vertices in a graph is equal to twice the no. of edges in the graph.

Given: A graph with some edges & vertices
e.g.



To prove: -

$$\sum \text{deg}(v) = 2e$$



proof :-

$$\begin{aligned} \sum \text{deg}(v) &= \text{deg}(v_1) + \text{deg}(v_2) + \text{deg}(v_3) \\ &\quad + \text{deg}(v_4) + \text{deg}(v_5) \\ &= 4 + 4 + 4 + 3 + 1 \end{aligned}$$

$$\sum \text{deg}(v) = 16 \quad \text{--- (i)}$$

$$\text{No. of edges} = 8 \quad \text{--- (ii)}$$

from (i) & (ii)

$$\sum \text{deg}(v) = 2e \quad \text{is proved}$$

Theorem:- In any graph, the no. of vertices of odd degree is always even. RGIPV Jan + Jun 2006/02/05

proof:- Let $G(V, E)$ be graph with

V : Set of vertices in G

E : Set of Edges in G

Let set of vertices of Even degree be V_e and that of odd degree be V_o

$$V_e \cup V_o = V$$

$$V_e \cap V_o = \phi$$

Hence

$$\sum_{v \in V} \text{deg}(v) = \sum_{v \in V_e} \text{deg}(v) + \sum_{v \in V_o} \text{deg}(v) \quad \text{--- (i)}$$

$$\therefore \sum_{v \in V_e} \text{deg}(v) = \text{even} = 2k \quad \text{--- (ii)}$$

$$\therefore \text{we know that } \sum \text{deg}(v) = 2e \quad \text{--- (iii)}$$



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Isom (i), (ii) & (iii)
we have

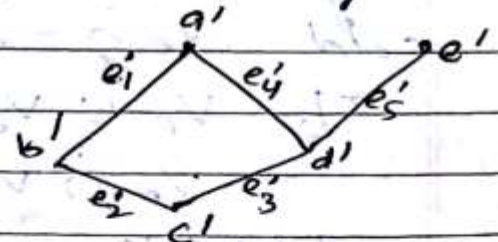
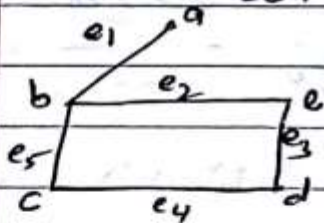
$$2e = 2k + \sum_{v \in V_0} \text{deg}(v)$$

$$\sum_{v \in V_0} \text{deg}(v) = 2(e-k) = \text{an Even no.}$$

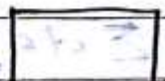
Hence the no. of terms must be even i.e. the no. of vertices of odd degree is always even.

Isomorphic: If two graphs have same Graph theoretic properties i.e.

- (i) They have same no. of edges
- (ii) They have same no. of vertices
- (iii) They have equal vertices with the given degree



Finite & Infinite Graphs: A graph with finite no. of edges and vertices is called a finite graph otherwise it is an infinite graph.



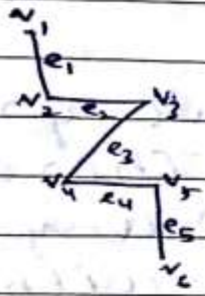
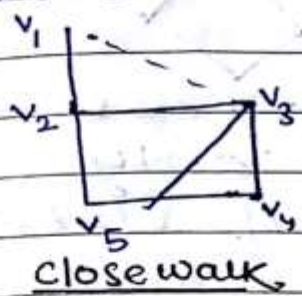
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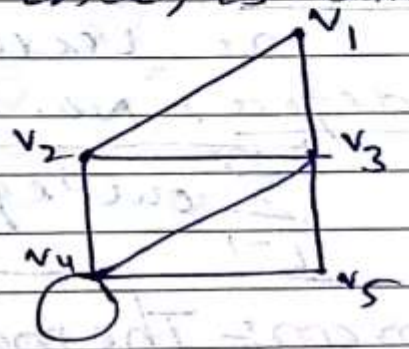
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walk:- A walk of graph 'G' is a finite sequence of vertices & edges, beginning & ending with same or diff. vertices, such that each edge is incident with the vertices preceding it. Such that no edge appears more than once. A vertex however may appear more than once.



If terminated vertex are diff. open walk

path:- An open walk in which no vertex appears more than once, is called path.



circuit:- A closed walk in which no (cycle) vertex appears more than once is called circuit. (path) C_k starts and ends with same vertex



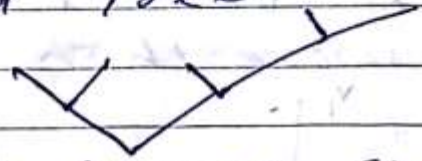
Euler Graph:- If a close walk in a graph contains all the edges of



the graph then the walk is called Euler line & such graph is called Euler graph.



Trees:- A connected graph without any circuit is called Tree



For finite graph trees are finite prop. ① There is only one path b/w every pair of vertices
② Tree with n vertices has $(n-1)$ edges.

In Degree:- No. of edges terminating at vertex.

out degree:- No. of edges beg. at vertex.

$$\sum_{i=1}^n \text{outdegree}(v_i) = \sum_{i=1}^n \text{indegree}(v_i)$$

Theorem:- The maximum no. of edges with n vertices in a simple graph is $\frac{n(n-1)}{2}$

Proof:- Let $G(V, E)$ has n vertices v_1, v_2, \dots, v_n
The vertex v_1 can be joined to $n-1$ vert v_2, v_3, \dots, v_n
similarly v_2 can be joined to $(n-2)$ vertices
similarly v_{n-1} ——— to only one new edge (v_{n-1}, v_n)

$$\text{Max. Edge} = (n-1)(n-2) + \dots + 2 + 1 = \frac{1}{2} n(n-1)$$



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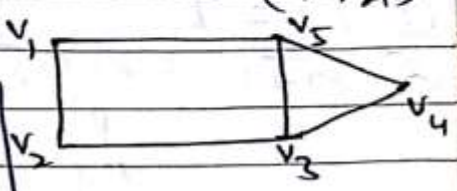
Matrix Representation of Graph

(a) Adjacency Matrix $a_{ij} = \begin{cases} 1 & \text{if there is an edge b/w } i^{\text{th}} \text{ \& } j^{\text{th}} \text{ vertices} \\ 0 & \text{if there is no edge} \end{cases}$

Q:- Find Adjacency Matrix (MA) of graph G

Soln:-

$$M_A = \begin{matrix} & \begin{matrix} v_1 & v_2 & v_3 & v_4 & v_5 \end{matrix} \\ \begin{matrix} v_1 \\ v_2 \\ v_3 \\ v_4 \\ v_5 \end{matrix} & \begin{bmatrix} 0 & 1 & 0 & 0 & 1 \\ 1 & 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 1 & 0 \end{bmatrix} \end{matrix}$$

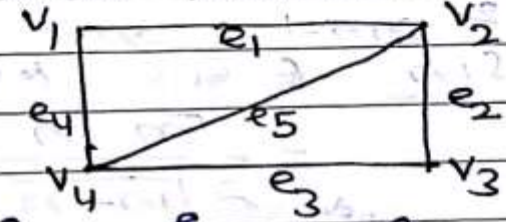


(b) Incidence Matrix :- If $G(V, E)$ then Incidence Matrix is $[A]_{n \times m} = [a_{ij}] = \begin{cases} 1 & \text{j}^{\text{th}} \text{ edge incident on } i^{\text{th}} \text{ vertex} \\ 0 & \text{otherwise} \end{cases}$

Q:- Find the Incidence Matrix 'A' of graph G

Soln

$$A = \begin{matrix} & \begin{matrix} e_1 & e_2 & e_3 & e_4 & e_5 \end{matrix} \\ \begin{matrix} v_1 \\ v_2 \\ v_3 \\ v_4 \end{matrix} & \begin{bmatrix} 1 & 0 & 0 & 1 & 0 \\ 1 & 1 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 & 1 \end{bmatrix} \end{matrix}$$



Note:- For Directed graph, if dirⁿ is towards the vertex consider -1 & similar to earlier case